第10章 An Example - The Countdown Problem

♦ What Is Countdown?

A popular quiz programme on British television that has been running since 1982.

Based upon an original French version called "Des Chiffres et Des Lettres".

Includes a numbers game that we shall refer to as the countdown problem.

◆ An example of Countdown

```
Using the numbers: 1 3 7 10 25 50 and the arithmetic operators: + - * ÷ construct an expression whose value is 765
```

◆ Two rules of Countdown?

- 1. All the numbers, including intermediate results, must be positive naturals (1, 2, 3, ...).
- 2. Each of the source numbers can be used at most once when constructing the expression.
- The example of Countdown: 1 3 7 10 25 50 ⇒ 765 One possible solution is: (25 - 10) * (50 + 1) = 765 There are 780 solutions for this example. Changing the target number to 831 gives an example that has no solutions.
- ◆ Evaluating Expressions

A type for Operators:

```
data Op = Add | Sub | Mul | Div deriving (Show)
```

Apply an operator:

```
apply :: 0p \rightarrow Int \rightarrow Int \rightarrow Int

apply Add x y = x + y

apply Sub x y = x - y

apply Mul x y = x * y

apply Div x y = x `div` y
```

Decide if the result of applying an operator to two positive natural numbers is another such:

```
valid :: Op → Int → Int → Bool
valid Add _ _ = True
valid Sub x y = x > y
valid Mul _ _ = True
valid Div x y = x `mod` y = 0
```

A type for Expressions:

```
data Expr = Val Int | App Op Expr Expr deriving (Show)
```

Return the overall value of an expression, provided that it is a positive natural number:

```
eval :: Expr \rightarrow [Int] eval (Val n) = [ n | n > 0 ] eval (App o l r) = [ apply o x y | x \leftarrow \text{eval l}, y \leftarrow \text{eval r}, \text{valid o x y}]
```

- either: succeeds and returns a singleton list.
- or: fails and returns the empty list.
- ◆ Some combinatorial functions

Returns all subsequences of a list.

```
subs :: [a] \rightarrow [[a]]

subs [] = [[]]

subs (x:xs) = let yss = subs xs in yss ++ map (x:) yss

ghci> subs [1, 2, 3]

[ [], [3], [2], [2, 3], [1], [1, 3], [1, 2], [1, 2, 3] ]
```

Returns all possible ways of inserting a new element into a list.

```
interleave :: a \rightarrow [a] \rightarrow [[a]]

interleave x [] = [[x]]

interleave x (y:ys) = (x:y:ys) : map (y:) (interleave x ys)
```

ghci> interleave 1 [2, 3, 4]

```
[[1,2,3,4], [2,1,3,4], [2,3,1,4], [2,3,4,1]]
```

Returns all permutations of a list.

```
perms :: [a] \rightarrow [[a]]

perms [] = [[]]

perms (x:xs) = concat $ map (interleave x) (perms xs)

ghci> perms [1, 2, 3]

[ [1,2,3], [2,1,3], [2,3,1], [1,3,2], [3,1,2], [3,2,1] ]
```

Return a list of all possible ways of choosing zero or more elements from a list in any order.

```
choices :: [a] \rightarrow [[a]] choices = concat . map perms . subs ghci> choices [1, 2, 3] [ [], [3], [2], [2,3], [3,2], [1], [1,3], [3,1], [1,2], [2,1], [1,2,3], [2,1,3], [2,3,1], [1,3,2], [3,1,2], [3,2,1] ]
```

→ Formalizing the countdown problem

Return a list of all the values in an expression.

```
values :: Expr \rightarrow [Int]
values (Val n) = [n]
values (App _ l r) = values l \leftrightarrow values
```

Decide if an expression is a solution for a countdown problem (i.e., a given list of source numbers, and a target number).

```
solution :: Expr \rightarrow [Int] \rightarrow Int \rightarrow Bool solution e ns n = (values e) `elem` (choices ns) && eval e = [n]
```

◆ Brute Force Solution

Return a list of all possible ways of splitting a list into two non-empty parts.

```
split :: [a] → [([a],[a])]
split [] = []
split [_] = []
split (x:xs) = ([x],xs) : [ (x:ls, rs) | (ls,rs) ← split xs ]
ghci> split [1, 2, 3, 4]
```

```
ghci> split [1, 2, 3, 4]
[ ([1], [2,3,4]), ([1,2], [3,4]), ([1,2,3], [4]) ]
```

Return a list of all possible expressions whose values are precisely a given list of numbers.

Return a list of all possible expressions that solve an instance of the countdown problem.

```
solutions :: [Int] \rightarrow Int \rightarrow [Expr] solutions ns n = [ e | ns' \leftarrow choices ns , e \leftarrow exprs ns' , eval e = [n] ]
```

♦ How Fast Is It?

```
Hardware: 2.8GHz Core 2 Duo, 4GB RAM
Compiler: GHC version 7.10.2
Example: solutions [1,3,7,10,25,50] 765
One solution: 0.108 seconds
All solutions: 12.224 seconds
(如果运行在ghci中,时间估计会增加一个数量级)
```

◆ Can We Do Better?

Many of the expressions that are considered will typically be invalid - fail to evaluate.

• For our example, only around 5 million of the 33 million possible expressions are valid.

Combining generation with evaluation would allow earlier rejection of invalid expressions.

→ Fusing generation and evaluation

A type for Valid expressions and their values:

```
type Result = (Expr, Int)
```

A function without fusion:

```
results :: [Int] \rightarrow [Result] results ns = [ (e,n) | e \leftarrow exprs ns, n \leftarrow eval e ]
```

A function with fusion:

♦ A better solution

```
solutions' :: [Int] \rightarrow Int \rightarrow [Expr] solutions' ns n = [ e | ns' \leftarrow choices ns , (e,m) \leftarrow results ns', m = n ]
```

♦ How Fast Now ?

Hardware: 2.8GHz Core 2 Duo, 4GB RAM

Compiler: GHC version 7.10.2

Example: solutions [1,3,7,10,25,50] 765

One solution: 0.108 s 0.014 s

All solutions: 12.224 s 1.312 s

Brute Force Fusion

◆ Can We Do Better Further ?

Many expressions will be essentially the same using simple arithmetic properties, such as:

```
; x * y = y * x
; x * 1 = x
```

Exploiting such properties would considerably reduce the search and solution spaces.

◆ A better valid function

```
valid :: Op → Int → Int → Bool
valid Add x y = x ≤ y
valid Sub x y = x > y
valid Mul x y = x ≤ y && x ≠ 1 && y ≠ 1
valid Div x y = x `mod` y = 0 && y ≠ 1

-- old one
valid :: Op → Int → Int → Bool
valid Add x y = True
valid Sub x y = x > y
valid Mul x y = True
valid Div x y = x `mod` y = 0
```

♦ How Fast Now ?

Hardware: 2.8GHz Core 2 Duo, 4GB RAM
Compiler: GHC version 7.10.2

Example:	solutions [1,3,7,10,25,50] 765		
One solution:	0.108 s	0.014 s	0.007 s
All solutions:	12.224 s	1.312 s	0.119 s
	Brute Force	Fusion	better valid

作业01

Modify the final program to:

- 1. allow the use of exponentiation in expressions;
- 2. produce the nearest solutions if no exact solution is possible;
- 3. order the solutions using a suitable measure of simplicity.